

Formal Logic

Lecture 9: The Semantics of Predicate Logic (Part II)

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Validity and other Notions in L_2

Logical truth in L_2

- A sentence ϕ in L_2 is *logically true* IFF it is true under any interpretation.

NB: Bergmann et al. (2014) call it ‘quantificationally true’.

Examples:

$$(\forall x)(Px \rightarrow Px)$$

$$(\exists x)(Px \vee \neg Px)$$

- We cannot show that a sentence is logically true via one or even some interpretations. We need all of them.
- We can show that a sentence is *not* logically true by coming up with one interpretation that makes the sentence false.

Contradiction in L_2

- A sentence ϕ in L_2 is *logically false* IFF it is false under any interpretation.

NB: Bergmann et al. (2014) call it ‘quantificationally false’.

Examples:

$$\neg(\exists x)(Px \vee \neg Px)$$

$$(\forall x) Px \ \& \ \neg Pa$$

- We cannot show that a sentence is logically false via one or even some interpretations. We need all of them.
- We can show that a sentence is *not* logically false by coming up with one interpretation that makes the sentence true.

Contingency in L_2

- A sentence ϕ in L_2 is *contingent* if and only if it is neither logically true nor logically false.

NB: Bergmann et al. (2014) call it ‘quant. indeterminate’.

Examples:

$$(\exists x)(Px \ \& \ Qx)$$

$$(\forall y) \text{ Ray}$$

- We can show that a sentence is contingent via two interpretations (one where it comes out true and one false).

Validity in L_2

- An argument in L_2 is *valid* IFF there is no interpretation under which the premises are true and the conclusion is false.

NB: Bergmann et al. (2014) call it ‘quantificationally valid’.

Examples:

$(\forall x) Px$	entails	Pa
$\{(\forall x)(Px \rightarrow Qa), (\exists x)Px\}$	entails	Qa

- We cannot show that an argument is valid via one or even some interpretations. We need all of them.
- We can show that an argument is invalid via 1 interpretation that makes the premises true and the conclusion false.

Consistency in L_2

- A set of sentences in L_2 is *consistent* IFF there is at least one interpretation on which all the members of the set are true.

NB: Bergmann et al. (2014) call them ‘quant. consistent’.

Examples:

$$\{(\exists x)Hx, \neg Hc\}$$

$$\{(\forall y)Ray, \neg Rba \vee (\exists z) \neg Raz\}$$

- We cannot show that a set of sentences is inconsistent via one or even some interpretations. We need all of them.
- We can show that a set of sentences is consistent via one interpretation that makes all sentences true.

Logical equivalence in L_2

- Two sentences ϕ, ψ in L_2 are *logically equivalent* IFF there is no interpretation on which ϕ, ψ have different truth-values.

NB: Bergmann et al. (2014) call them ‘quant. equivalent’.

Examples:

$$\begin{array}{lll} (\forall x) \neg Qx & \text{and} & \neg(\exists x) Qx \\ (\exists x)(Ax \ \& \ \neg Bx) & \text{and} & \neg(\forall x)(Ax \rightarrow Bx) \end{array}$$

- We cannot show that a pair of sentences is equivalent via one or even some interpretations. We need all of them.
- We can show that a pair is *not* logically equivalent via 1 interpretation where one sentence is true and the other false.

The undecidability of L_2

- Alonzo Church (1936) proved a fundamental limit to first-order predicate logic.
 - There is no decision procedure that determines for every group of L_2 sentences whether they are:
 - * logically true, logically false or contingent
 - * consistent or inconsistent
 - * valid or invalid
 - * equivalent or inequivalent
- NB:** In L_1 , there is such a procedure, namely truth-tables.
- **Decision procedure:** A mechanical method that shows for all cases & in a finite number of steps whether a property holds.

Exercise Set #5

Exercise 4.1

- **Instructions:** Determine whether the following expressions are formulae of L_2 and say which of those are also sentences of L_2 .

Exercise 4.1

Formula, sentence or neither?

- (i) $(\forall x) (P_1x \rightarrow Qy)$
- (ii) $(\exists x) \neg(\neg(\exists y)Py \ \& \ \neg\neg\neg Rxa)$
- (iii) P
- (iv) $(\forall x)(\exists y)(\exists z) (R_{25}xyz)$
- (v) $(\forall x)(\exists x) Qxx$
- (vi) $\neg(\neg((\exists x)Px \ \& \ (\exists y)Qy))$
- (vii) $(\forall x)(\exists y(Pxy \ \& \ Px) \vee Qxyx)$

Exercise 4.1: Solution

Formula, sentence or neither?

- | | |
|--------------------------------------------------------------------|-------------------------------------------|
| (i) $(\forall x) (P_1x \rightarrow Qy)$ | Formula (y is free) |
| (ii) $(\exists x) \neg(\neg(\exists y)Py \ \& \ \neg\neg\neg Rxa)$ | Sentence |
| (iii) P | Sentence |
| (iv) $(\forall x)(\exists y)(\exists z) (R_{25}xyz)$ | Neither (brackets on R) |
| (v) $(\forall x)(\exists x) Qxx$ | Sentence |
| (vi) $\neg(\neg((\exists x)Px \ \& \ (\exists y)Qy))$ | Neither (outer brackets) |
| (vii) $(\forall x)((\exists y)(Pxy \ \& \ Px) \vee Qxyx)$ | Formula (2nd y is free) |

Exercise 4.3

- **Instructions:** Provide L_2 -formalisations for the following English sentences.

Exercise 4.3

English Sentences:

- (i) London is big and ugly.
 - (ii) Culham is a large village.
 - (iii) A city has a city hall.
 - (iv) Material objects are divisible.
 - (v) Tom owns at least one car.
 - (vi) Tom owns at least one car and
he won't sell it.
 - (vii) One man has visited every country.

Exercise 4.3: Solution

English Sentences:

- (i) London is big and ugly.
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Formalised:

- Ba&Ua
- Lc
- $(\forall x)(Cx \rightarrow (\exists y)(Pxy \& Hy))$
- $(\forall x)(Mx \rightarrow Dx)$
- $(\exists x)(Otx \& Cx)$
- $(\exists x)((Otx \& Cx) \& \neg Stx)$
- $(\exists x)(Mx \& (\forall y)(Cy \rightarrow Vxy))$

NB: Other readings sometimes possible, e.g. one means one in vii.

Exercise 4.4

- **Instructions:** Translate the L_2 -sentences below into English using the following dictionary.

Exercise 4.4

- Translate the L_2 -sentences below into English using the following dictionary.

a: Tom

P¹: ... is a person

Q¹: ... acts freely

L_2 Sentences:

- (i) Qa
- (ii) (Qa \vee \neg Pa)
- (iii) (\forall x) (Px \rightarrow Qx)
- (iv) (\forall x) (Px \leftrightarrow Qx)
- (v) \neg \exists z₁Qz₁

Exercise 4.4: Solution

- Translate the L_2 -sentences below into English using the following dictionary.

a: Tom

P¹: ... is a person

Q¹: ... acts freely

L_2 Sentences:

(i) Qa

(ii) (Qa \vee \neg Pa)

(iii) (\forall x) (Px \rightarrow Qx)

(iv) (\forall x) (Px \leftrightarrow Qx)

(v) \neg \exists z₁Qz₁

English Sentences:

Tom acts freely.

Tom acts freely or Tom is not a person.

Every person acts freely.

Every person acts freely and vice-versa.

Nothing acts freely.

Exercise 5.1

- **Instructions:** Consider an L_2 -structure S ... Are the following sentences true or false in this structure? Sketch proofs...

Exercise 5.1

- Consider an L_2 -structure S ... Are the following sentences true or false in this structure? Sketch proofs...

UD: $\{1, 2, 3\}$, P: $\{2\}$, R: $\{\langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 1, 3 \rangle\}$, a: 1, b: 3

- (i) Pa
- (ii) Rab
- (iii) Rba
- (iv) $Rab \leftrightarrow Rba$
- (v) $Rbb \vee (\neg Pa \wedge \neg Raa)$
- (vi) $(\exists x)Rax$
- (vii) $(\exists x)(Rax \wedge Rx b)$
- (viii) $Pb \vee (\exists x)Rxx$
- (ix) $(\forall x)(\exists y)Rxy$

Exercise 5.1: Solution

- Consider an L_2 -structure S ... Are the following sentences true or false in this structure? Sketch proofs...

UD: $\{1, 2, 3\}$, P: $\{2\}$, R: $\{\langle 1, 2 \rangle, \langle 2, 3 \rangle, \langle 1, 3 \rangle\}$, a: 1, b: 3

(i) Pa	False as $1 \notin P$
(ii) Rab	True as $\langle 1, 3 \rangle \in R$
(iii) Rba	False as $\langle 3, 1 \rangle \notin R$
(iv) $Rab \leftrightarrow Rba$	False as Rab is true and Rba false
(v) $Rbb \vee (\neg Pa \wedge \neg Raa)$	True as $1 \notin P$ and $\langle 1, 1 \rangle \notin R$
(vi) $(\exists x)Rax$	True, e.g. $\langle 1, 2 \rangle \in R$
(vii) $(\exists x)(Rax \wedge Rx b)$	True as $\langle 1, 2 \rangle \in R$ and $\langle 2, 3 \rangle \in R$
(viii) $Pb \vee (\exists x)Rxx$	False as neither $3 \in P$ nor $\langle x, x \rangle \in R$
(ix) $(\forall x)(\exists y)Rxy$	False as $\langle 3, y \rangle \notin R$

The End